# Seaform: Search-As-You-Type in Forms

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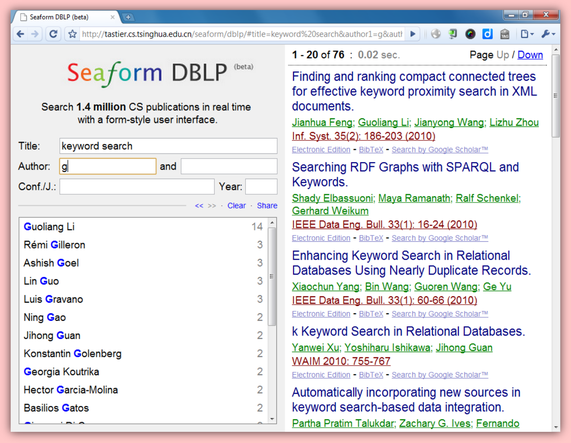
## What is Seaform?

Seaform is a keyword search system that can search a relational table in real time using a form-style user interface.

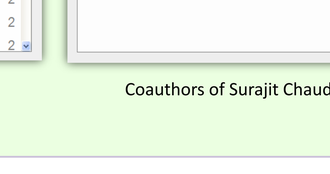



## Why Seaform?

#1: SQL is complex.

```
SELECT *
FROM   Author A, Autor_Paper AP, Paper P
WHERE  title   LIKE 'keyword' AND
       title   LIKE 'search'  AND
       authors LIKE 'g%'      AND
       A.id    =    AP.aid    AND
       P.id    =    AP.pid
```

#2: Form is awkward.




#3: Traditional keyword search is imprecise.

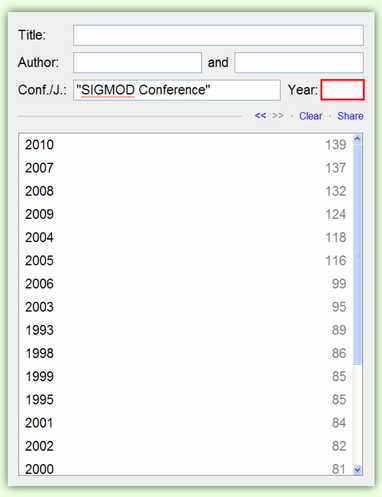



## Features

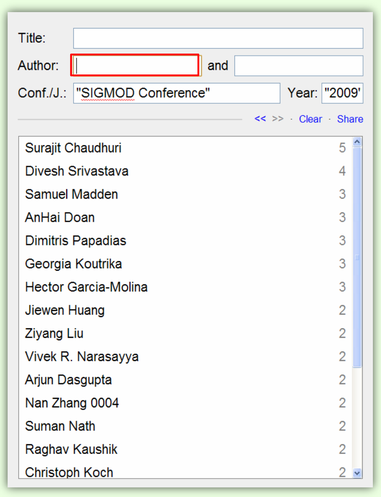
#1: Real-time response: less than 30 ms. for each keystroke in average.

#2: Precise search conditions.

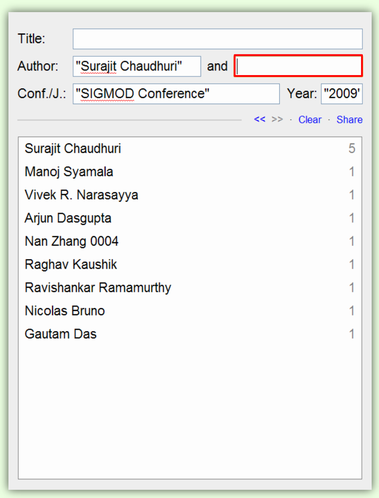
#3: Faceted search.




# of accepted papers over years.




All authors sorted by # of papers.




Coauthors of Surajit Chaudhuri.

## System Overview

Data:

| ID | Title | Conf. | Author |
|---|---|---|---|
| 1 | xml database | VLDB | albert |
| 2 | xml database | SIGMOD | bob |
| 3 | xml search | VLDB | albert |
| 4 | xml security | VLDB | alice |
| 5 | rdbms | SIGMOD | charlie |

Global table

| ID | Title |
|---|---|
| T1 | xml database |
| T2 | xml search |
| T3 | xml security |
| T4 | rdbms |

Local tables

| ID | Conf. |
|---|---|
| C1 | VLDB |
| C2 | SIGMOD |

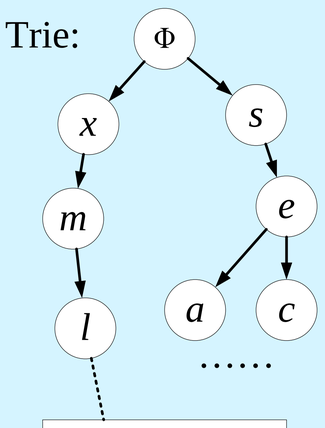
| ID | Author |
|---|---|
| A1 | albert |
| A2 | bob |
| A3 | alice |
| A4 | charlie |

### Index:

For each local table, we build:

1. A trie structure;
2. A local-global mapping table; and
3. A global-local mapping table.

Trie:




L-G Mapping Table:

| | |
|---|---|
| T1 | 1, 2 |
| T2 | 3 |
| T3 | 4 |
| T4 | 5 |

G-L Mapping Table:

| | |
|---|---|
| 1 | T1 |
| 2 | T1 |
| 3 | T2 |
| 4 | T3 |
| 5 | T4 |

### Search Algorithm:

1. Identify $\Delta q = q - q_0$ using cache;
2. Calculate $L_q^* = \text{TrieSearch}(\Delta q) \cap L_{q_0}$;
3. Calculate $G_q = \text{MapToGlobal}(L_q^*) \cap G_{q_0}$;
4. Calculate $L_q = \text{MapToLocal}(G_q) \cap L_{q_0}$.

*Synchronization*

T1: xml database
T2: xml search
T3: xml security

Title: xml
Author:
xml database
xml search
xml security
xml database (albert)
xml database (bob)
xml search (albert)
xml security (alice)

$\Delta q$ = title:xml
$L_q^*$ = {xml database, xml search, xml security}
$G_q$ = {xml database (albert), xml database (bob), xml search (albert), xml security (alice)}
$L_q$ = {xml database, xml search, xml security}

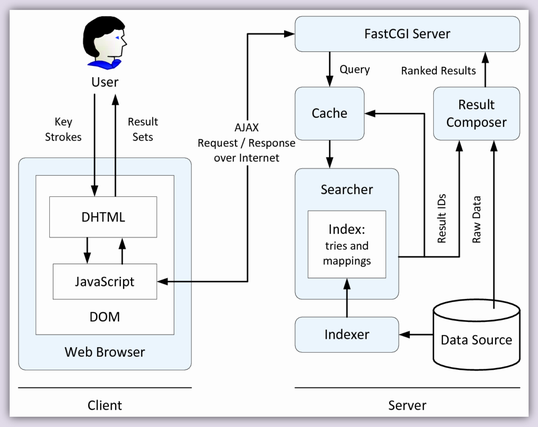
Title: xml
Author:
albert
bob
alice
xml database (albert)
xml database (bob)
xml search (albert)
xml security (alice)

$\Delta q$ = author:
$L_q^*$ = {albert, bob, alice, charlie}
$G_q$ = {xml database (albert), xml database (bob), xml search (albert), xml security (alice)}
$L_q$ = {albert, bob, alice}

Title: xml
Author: al
albert
alice
xml database (albert)
xml search (albert)
xml security (alice)

$\Delta q$ = author:al
$L_q^*$ = {albert, alice}
$G_q$ = {xml database (albert), xml search (albert), xml security (alice)}
$L_q$ = {albert, alice}

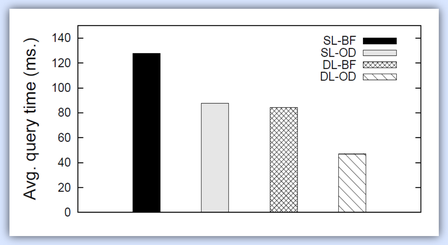
## Implementation




- Client/server architecture:
  1) Client: Javascript/Ajax,
  2) Server: FastCGI.
- Datasets:
  1) DBLP (1.4m publications),
  2) IMDB (507k movies).

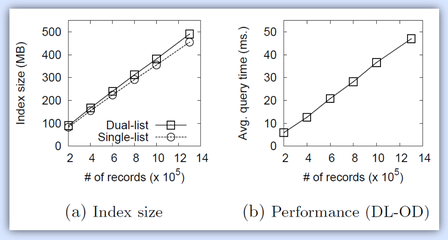
## Experiments

Comparison of four algorithms:




SL: single-list tire, DL: dual-list trie,
BF: brute-force sync., OD: on-demand sync.

Scalability of Seaform-DBLP:




(a) Index size (b) Performance (DL-OD)

http://tastier.cs.thu.edu.cn/seaform/

VLDB Demonstration Track, Singapore, Sept. 2010.